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The Solvability Complexity Index of Sampling-based Hilbert Transform Approximations

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Intern. Conf. on Sampling Theory and Applications July 11th, 2019 – Bordeaux, France



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Outline of this Talk

- 1. The Solvability Complexity Index and Towers of Algorithms
- 2. The Hilbert Transform and Sampling-based Hilbert Transform Approximations
- 3. Apply the the SCI Framework to the Hilbert Transform
- 4. Conlusions

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The Solvability Complexity Index

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The Solvability Complexity Index

- The Solvability Complexity Index (SCI) is a new measure to characterize the complexity of certain computational problems.
 - J. Ben-Artzi, A. C. Hansen, O. Nevanlinna, and M. Seidel, "Can everything be computed? On the Solvability Complexity Index and towers of algorithms." *preprint: arXiv:1508.03280v1*, Aug. 2015.
 - J. Ben-Artzi, A. C. Hansen, O. Nevanlinna, and M. Seidel, "New barriers in complexity theory: On the solvability complexity index and the towers of algorithms," *Comptes Rendus Mathematique*, vol. 353, no. 10, pp. 931–936, 2015.
- ▷ Solving computational problems by towers of algorithms.
- \triangleright The SCI is the minimal height of such a tower.
- ▷ The SCI counts the minimum number of limits necessary for computing a desired quantity.



The Solvability Complexity Index – Polynomials

- ▷ Let \mathscr{P}_N set of polynomials of degree $N \in \mathbb{N}$.
- ▷ Can we find an algorithm to compute the roots of every $p \in \mathscr{P}_N$ using only finitely many arithmetic operations?
 - Yes! for $N \leq 4$
 - No! for $N \ge 5$
- ▷ What if we allow for a limiting process? Does there exists a sequence of operations $\{\Gamma_n : \mathscr{P}_N \to \mathbb{C}\}$ where each Γ_n involves only finitely many arithmetic operations such that

$$\lim_{n\to\infty}\Gamma_n(p)=z_0$$

is a zero of *p*?

− Yes! for all $N \in \mathbb{N}$

Note: Newton's method will not work.
 An there is no "purely iterative generally convergent algorithm" for root finding.

S. Smale, "The fundamental theorem of algebra and complexity theory," Bull. Amer. Math. Soc., vol. 4, 1981.

P. Doyle and C. McMullen, "Solving the quintic by iteration," *Acta Math.*, vol. 163, pp. 151–180, 1989.

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The Solvability Complexity Index

- b There are important computational problems which can not be solved by one limiting process!
 - "purely iterative generally convergent algorithm" for root finding.
 - ...
 - calculation of the Hilbert transform from samples.
- What if we allow for a several limiting processes?
 Does there exist a sequence {Γ_{n_k,...,n₁}} of operations depending on several indices n_k,...,n₁ ∈ N such that

$$\lim_{n_1\to\infty}\cdots\lim_{n_k\to\infty}\Gamma_{n_k,\dots,n_1}(p)=\Xi(p)$$

converges to the desired result $\Xi(p)$?

H. Boche and V. Pohl, "On the calculation of the Hilbert transform from interpolated data," *IEEE Trans. Inf. Theory*, vol. 54, no. 5, pp. 2358–2366, May 2008.

H. Boche and V. Pohl, "Calculating the Hilbert transform on spaces with energy concentration: Convergence and divergence regions," *IEEE Trans. Inf. Theory*, vol. 65, no. 1, pp. 586–603, Jan. 2019.
 Volker Pohl (TUM) | The Solvability Complexity Index of Sampling-based Hilbert Transform Approximations| SampTA 2019

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The SCI Framework

The SCI framework, is based on the following basic objects

- Ω a primary set, is the set of objects on which we perform the computations
- ℳ a (pseudo) metric space
 - Ξ a problem function, is the mapping $\Xi: \Omega \to \mathscr{M}$ which we want to compute for our objects in Ω
 - A is the evaluation set, that is a set of (measurement) functionals $\lambda:\Omega\to\mathbb{C}$

Definition (Computational problem)

A collection $\{\Xi, \Omega, \mathcal{M}, \Lambda\}$ is said to be a *computational problem*.

Example ({ $\mathscr{C}(\mathbb{T}), \mathbb{R}, \|\cdot\|_{\infty}, \mathbb{T}$ } – "Sampling-based computational problem")

- $\Omega = \mathscr{C}(\mathbb{T})$, the Banach space of continuous functions on a compact interval $\mathbb{T} \subset \mathbb{R}$
- $\mathcal{M} = \mathbb{R}$ the real numbers.
- $\Xi : \mathscr{C}(\mathbb{T}) \to \mathbb{R}$ is given by $\Xi(f) = \|f\|_{\infty} = \max_{t \in \mathbb{T}} |f(t)|$
- $\Lambda = \{\lambda_\tau : \tau \in \mathbb{T}\}$ consists of all point evaluations

$$\lambda_{ au}(f)=f(au)$$
 for all $f\in \mathscr{C}(\mathbb{T})$.

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Fundamental Algorithm

- \triangleright Let $\{\Xi, \Omega, \mathcal{M}, \Lambda\}$ be a given computational problem.
- ▷ The problem of calculating $\Xi(f)$ will be done by a so called tower of algorithms.
- ▷ The next definition characterizes the algorithms at the lowest (the fundamental) level of the tower.

Definition (Fundamental Algorithm)

Let $\{\Xi, \Omega, \mathscr{M}, \Lambda\}$ be a computational problem. A mapping $\Gamma : \Omega \to \mathscr{M}$ is said to be a *fundamental algorithm* if for every $f \in \Omega$ the following conditions are satisfied:

- (i) There exists a finite subset $\Lambda_{\Gamma}(f) \subset \Lambda$ of evaluations.
- (ii) The action of Γ on f depends only on the values $\{\lambda(f) : \lambda \in \Lambda_{\Gamma}(f)\}$.
- (iii) For every $g \in \Omega$, with $\lambda(g) = \lambda(f)$ for all $\lambda \in \Lambda_{\Gamma}(f)$, holds $\Lambda_{\Gamma}(g) = \Lambda_{\Gamma}(f)$.

We say that Γ is a fundamental algorithm of type

- L for "*linear*", if Γ is linear.
- G for "general", if there is no further restriction on Γ .

Remark:

- The calculation of $\Gamma(f)$ is based on finitely many measurements of f.
- The evaluation set $\lambda_{\Gamma}(f)$ may be chosen adaptively, dependent on the actual object *f*.
- Condition (iii): well definiteness and consistence. It limits adaptivity.



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Tower of Algorithms

Definition (Tower of Algorithms)

A *tower of algorithms* of type α and of height k for $\{\Xi, \Omega, \mathcal{M}, \Lambda\}$ is a family of sequences of operators

$$\begin{split} & \Gamma_{n_k} : \Omega \to \mathcal{M} , \\ & \Gamma_{n_k,n_{k-1}} : \Omega \to \mathcal{M} , \\ & \vdots \\ & \Gamma_{n_k,n_{k-1},\dots,n_1} : \Omega \to \mathcal{M} , \end{split}$$

where $n_k, \ldots, n_1 \in \mathbb{N}$ and where the operators $\Gamma_{n_k, n_{k-1}, \ldots, n_1}$ at the lowest level of the tower are fundamental algorithms of type α . Moreover, for every $f \in \Omega$ holds

$$\Xi(f) = \lim_{n_k \to \infty} \Gamma_{n_k}(f)$$

$$\Gamma_{n_k}(f) = \lim_{n_{k-1} \to \infty} \Gamma_{n_k, n_{k-1}}(f)$$

$$\vdots$$

$$\Gamma_{n_k, \dots, n_2}(f) = \lim_{n_1 \to \infty} \Gamma_{n_k, \dots, n_2, n_1}(f)$$

where the convergence is always in the metric of \mathcal{M} .



Solvability Complexity Index

Definition (Solvability Complexity Index)

We say that a computational problem $\{\Xi, \Omega, \mathcal{M}, \Lambda\}$ has a *Solvability Complexity Index*

 $SCI(\Xi,\Omega,\mathcal{M},\Lambda)_{\alpha} = k$

with respect to towers of algorithms of type α if *k* is the smallest integer for which there exists a tower of algorithms of type α of height *k*. If no such tower exists then SCI $(\Xi, \Omega, \mathcal{M}, \Lambda)_{\alpha} = \infty$.

▷ We say that $SCI(\Xi, \Omega, \mathcal{M}, \Lambda)_{\alpha} = 0$, if $\Xi(f)$ can be calculate by finitely many algebraic operations.



Hilbert Transformation



The Hilbert Transform

Technical University of Munich

We consider continuous functions $f \in \mathscr{C}(\mathbb{T})$ on the interval $\mathbb{T} = [-\pi, \pi]$ with $f(-\pi) = f(\pi)$ \triangleright

The Fourier coefficients of $f \in \mathscr{C}(\mathbb{T})$ are given by \triangleright

$$c_k(f) = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(t) e^{-ikt} dt$$
, $k = 0, \pm 1, \pm 2, ...$

 \triangleright Assume that the Fourier series of f converges to f (in some sense)

$$f(t) = \sum_{k=-\infty}^{\infty} c_k(f) e^{ikt}, \qquad t \in \mathbb{T}.$$

The conjugate function associated with f is \triangleright

$$\widetilde{f}(t) = (\mathrm{H}f)(t) = -\mathrm{i}\sum_{k=-\infty}^{\infty} \mathrm{sgn}(k) c_k(f) \mathrm{e}^{\mathrm{i}kt}, \quad t \in \mathbb{T} \quad \text{with} \quad \mathrm{sgn}(k) = \begin{cases} 0 & : \ k = 0 \\ \frac{k}{|k|} & : \ k \neq 0 \end{cases}$$

0.4

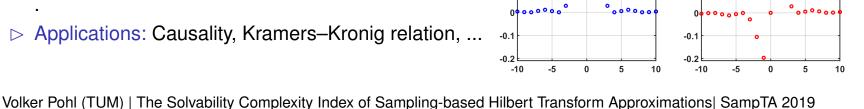
0.3

0.2

0.1

The Hilbert transform is the mapping \triangleright $H: f \mapsto \tilde{f}$





Fourier coefficients $c_n(f)$

Fourier coefficients $c_n(i\tilde{f})$

0.4

0.3

0.2

0.1



Signal Spaces for the Hilbert Transform

▷ Let $f \in \mathscr{C}(\mathbb{T})$ with Fourier coefficients $\{c_k(f)\}_{k \in \mathbb{Z}}$.

 \triangleright For $\beta \geq$ 0, we define the seminorm

$$\|f\|_{\beta} = \left(\sum_{k\in\mathbb{Z},\ k\neq 0} |k| (1+\log|k|)^{\beta} |c_k(f)|^2\right)^{1/2}.$$

 $\,\triangleright\,$ Therewith, we define for any $\beta\geq 0$ the Banach space

$$\mathscr{B}_{eta} = \left\{ f \in \mathscr{C}(\mathbb{T}) \ : \ \widetilde{f} \in \mathscr{C}(\mathbb{T}) \ \text{and} \ \|f\|_{eta} < \infty
ight\}$$

with the norm

$$\|f\|_{\mathscr{B}_{\beta}} = \max\left(\|f\|_{\mathscr{C}(\mathbb{T})}, \|\widetilde{f}\|_{\mathscr{C}(\mathbb{T})}, \|f\|_{\beta}\right) \,.$$

Remark

- For every $f \in \mathscr{B}_{\beta}$ the Fourier series and the conjugate Fourier series converge uniformly to f and \tilde{f}
- The space $\mathscr{B}_{\beta=0}$ corresponds to the usual Sobolev space $H^{1/2}(\mathbb{T}) = W^{1/2,2}(\mathbb{T})$.
- The value $||f||_{\beta=0}^2$ corresponds to the Dirichlet energy of *f*.
- $\mathscr{B}_{\beta_2} \subset \mathscr{B}_{\beta_1} \subset \mathscr{B}_{\beta=0}$ for all $\beta_2 \geq \beta_1 \geq 0$.

No Sampling-based Hilbert Transform Approximation

Theorem

Let $\{\Gamma_n\}_{n\in\mathbb{N}}$ be a sequence of lower semicontinuous mappings $\Gamma_n : \mathscr{B}_\beta \to \mathscr{C}(\mathbb{T})$ such that to every $n \in \mathbb{N}$ there exists a finite sampling set $\mathscr{T}_n \subset \mathbb{T}$ such that

 $f_1(\tau) = f_2(\tau)$ for all $\tau \in \mathscr{T}_n$ implies $\Gamma_n(f_1) = \Gamma_n(f_2)$.

Then for every $0 \le \beta \le 1$ there exist functions $f \in \mathscr{B}_{\beta}$ such that

 $\limsup_{n\to\infty} \left\|\widetilde{f}-\Gamma_n(f)\right\|_{\mathscr{C}(\mathbb{T})}>0.$

▷ The operators $\Gamma_n : \mathscr{B}_\beta \to \mathscr{C}(\mathbb{T})$ can be non-linear.

- ▷ So the Hilbert transform on \mathscr{B}_{β} can not be calculated by a single limit (i.e. SCI > 1).
- ▷ So what is the SCI for calculating the Hilbert transform?

H. Boche and V. Pohl, "Investigations on the approximability and computability of the Hilbert transform with applications," *Appl. Comput. Harmon. Anal.*, Sep. 2018, in press.



The Solvability Complexity Index for Calculating the Hilbert Transform



The Computational Problem

Now we investigate the Hilbert transform calculation using the SCI farmework.

 \triangleright For $eta \in [0,1]$, we consider the computational problems

 $\{H, \mathscr{B}_{\beta}, \mathscr{C}(\mathbb{T}), \Lambda\}$ and $\{H, \mathscr{B}_{\beta}, \mathscr{B}_{\beta}, \Lambda\}$

i.e.

- − The primary set Ω is \mathscr{B}_{β} for some $\beta \in [0, 1]$.
- The metric space \mathscr{M} is either $\mathscr{C}(\mathbb{T})$ or \mathscr{B}_{β} .
- The problem function Ξ is the Hilbert transform $H: \mathscr{B}_{\beta} \to \mathscr{M}$ is the Hilbert transform.
- The evaluations set Λ contains the sampling functionals $\lambda_{\tau}: \mathscr{B}_{\beta} \to \mathbb{C}$ with $\tau \in \mathbb{T}$ defined by

$$\lambda_{ au}(f)=f(au)\,,\quad f\in\mathscr{B}_{eta}\;.$$

▷ So the evaluation set consist now of point evaluations (i.e. sampling).



Sampling-based Fundamental Algorithms

Definition (Sampling-based Fundamental Algorithm)

Let $\{H, \mathscr{B}_{\beta}, \mathscr{M}, \Lambda\}$ be a sampling-based computational problem for calculating the Hilbert transform. A lower semicontinuous mapping $\Gamma : \mathscr{B}_{\beta} \to \mathscr{M}$ is said to be a *sampling-based fundamental algorithm* if

- (i) There exists a finite subset ${\it T}_{\Gamma} \subset \mathbb{T}$
- (ii) For all $f, g \in \mathscr{B}_{\beta}$

 $f(au) = g(au) \qquad ext{ for all } au \in \mathcal{T}_{\Gamma}$ implies $[\Gamma(f)](t) = [\Gamma(g)](t) \qquad ext{ for all } t \in \mathbb{T}$.

Moreover, we say that Γ is of type

- 1. *linear (L)*, if Γ is linear.
- 2. general (G), if there is no further restriction on Γ .

Remark

- \triangleright Every fundamental algorithm Γ is concentrated on a finite sampling set Λ_{Γ} .
- \triangleright No adaptivity, i.e. the sampling sets Λ_{Γ} depend only on Γ but not on the actual function.



The SCI of the Hilbert Transform

Theorem

Let $\beta \in [0,1]$ be arbitrary, then we have

$$\mathrm{SCI}(\mathrm{H},\mathscr{B}_{\beta},\mathscr{B}_{\beta},\Lambda)_{L}=\mathrm{SCI}(\mathrm{H},\mathscr{B}_{\beta},\mathscr{C}(\mathbb{T}),\Lambda)_{G}=2.$$

Remark

▷ We also get

$$2 = \operatorname{SCI}(\operatorname{H}, \mathscr{B}_{\beta}, \mathscr{C}(\mathbb{T}), \Lambda)_{G} \leq \operatorname{SCI}(\operatorname{H}, \mathscr{B}_{\beta}, \mathscr{C}(\mathbb{T}), \Lambda)_{L} \leq \operatorname{SCI}(\operatorname{H}, \mathscr{B}_{\beta}, \mathscr{B}_{\beta}, \Lambda)_{L} = 2$$
$$2 = \operatorname{SCI}(\operatorname{H}, \mathscr{B}_{\beta}, \mathscr{C}(\mathbb{T}), \Lambda)_{G} \leq \operatorname{SCI}(\operatorname{H}, \mathscr{B}_{\beta}, \mathscr{B}_{\beta}, \Lambda)_{G} \leq \operatorname{SCI}(\operatorname{H}, \mathscr{B}_{\beta}, \mathscr{B}_{\beta}, \Lambda)_{L} = 2$$

Non-linear methods give no improvement over linear approximation methods (as far as the SCI is concerned).

Proof

- ▷ Our previous already implied that $SCI(H, \mathscr{B}_{\beta}, \mathscr{B}_{\beta}, \Lambda)_{G} > 1$.
- ▷ It remains to show that SCI $(H, \mathscr{B}_{\beta}, \mathscr{B}_{\beta}, \Lambda)_{L} \leq 2$.



Tower of Algorithms for the Hilbert Transform

$$(\mathrm{H}f)(t) = -\mathrm{i}\sum_{k=-\infty}^{\infty} \mathrm{sgn}(k) \, c_k(f) \, \mathrm{e}^{\mathrm{i}kt} \,, \quad t \in \mathbb{T} \qquad \text{with} \qquad c_k(f) = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(t) \, \mathrm{e}^{-\mathrm{i}kt} \, \mathrm{d}t \,.$$

▷ Let $n_1, n_2 \in \mathbb{N}$ be fixed. For $f \in \mathscr{B}_\beta$, we consider

$$c_k(f, n_1) = rac{1}{n_1} \sum_{l=0}^{n_1-1} f\left(l rac{2\pi}{n_1}\right) \mathrm{e}^{-\mathrm{i} rac{2\pi}{n_1} l k} \,, \qquad ext{for all } |k| \le n_2 \,.$$

 $c_k(f, n_1)$ is the approximation of the *k*th Fourier coefficient of *f* by a Riemann sum based on n_1 equally spaced sampling points and we have

$$\lim_{n_1\to\infty}c_k(f,n_1)=c_k(f) \qquad \text{for all } |k|\leq n_2 \ .$$

▷ Define the operators in the lower level of the tower by

$$(\Gamma_{n_2,n_1}f)(t) = -i\sum_{k=-n_2}^{n_2} \operatorname{sgn}(k) c_k(f,n_1) e^{ikt}, \qquad t \in \mathbb{T},$$

▷ Define the operators in the upper level of the tower by

$$(\Gamma_{n_2}f)(t) = -\mathrm{i}\sum_{k=-n_2}^{n_2}\mathrm{sgn}(k) c_k(f) \mathrm{e}^{\mathrm{i}kt}, \qquad t \in \mathbb{T}.$$



Discussion

▷ Tower of algorithm

$$(\Gamma_{n_2}f)(t) = -i\sum_{k=-n_2}^{n_2} \operatorname{sgn}(k) c_k(f) e^{ikt}$$

$$(\Gamma_{n_2,n_1}f)(t) = -i\sum_{k=-n_2}^{n_2} \operatorname{sgn}(k) c_k(f,n_1) e^{ikt} \quad \text{with} \quad c_k(f,n_1) = \frac{1}{n_1} \sum_{l=0}^{n_1-1} f\left(l\frac{2\pi}{n_1}\right) e^{-i\frac{2\pi}{n_1}lk}.$$

▷ The first layer calculates the Fourier coefficients $c_k(f)$ of f from the samples $\left\{f\left(I\frac{2\pi}{n_1}\right)\right\}$ of f by the first limit

$$\lim_{n_1\to\infty}c_k(f,n_1)=c_k(f).$$

▷ The second layer calculates $\tilde{f} = Hf$ by the conjugate Fourier series (second limit).

- \triangleright How to get a tower of algorithm of height 1?
 - \Rightarrow Other evaluation set Λ (i.e. other measurement functionals)

$$\lambda_k: f\mapsto c_k(f), \qquad k=0,\pm 1,\pm 2,\ldots$$

 \Rightarrow Smaller primary set $\Omega \subset \mathscr{B}_{\beta}$.



Conclusions





Conclusions and Discussion

- ▷ The SCI for calculating the Hilbert transform H*f* from samples of $f \in \mathscr{B}_{\beta}$ is equal to 2.
 - One needs 2 limits to calculate $\tilde{f} = H(f)$.
 - The SCI = 2 for linear and general towers of algorithms.
 - This statement holds primary sets \mathscr{B}_{eta} with $eta\in[0,1]$
- $\,\triangleright\,$ For eta> 1 we (probably) get

$$\mathrm{SCI}(\mathrm{H},\mathscr{B}_{\beta},\mathscr{B}_{\beta},\Lambda)_{L}=1$$
.

- ▷ Similar results for other sampling based problems:
 - Recovering functions from samples on the space of uniformly convergent Fourier series.
 - Solving the Dirichlet problem based on samples of the boundary function.
 - Calculating the Spectral factorization.