

Third-Party Tokens for QUIC Address Validation

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QUIC

QUIC [RFC 9000] enables fast connection establishment for HTTP/3 in 1 RTT.

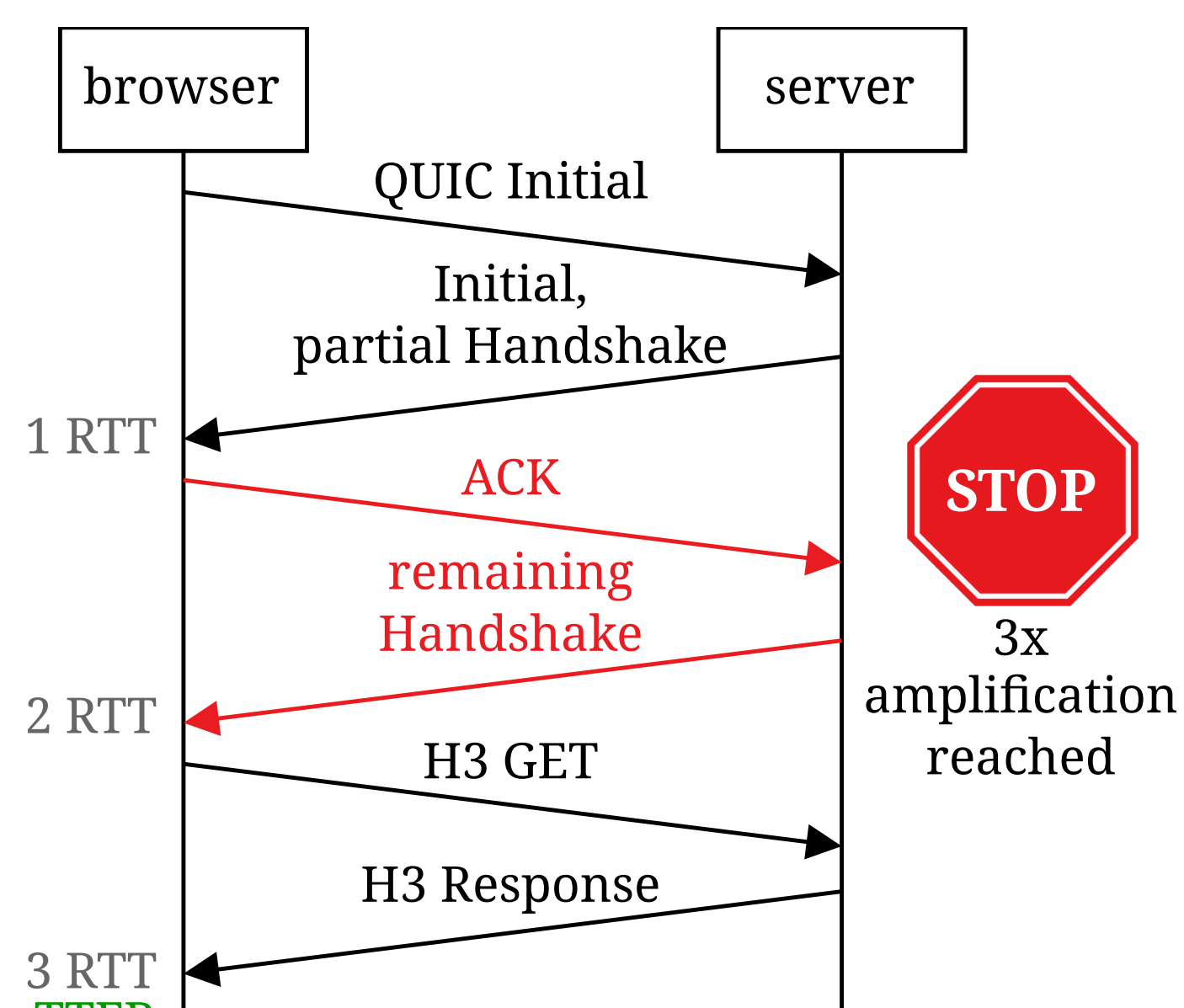
To prevent amplification attacks a QUIC server must limit the amount of data sent to an unvalidated address to 3x the amount of data received from the client.

The client address is validated after 1 RTT or with an Address Validation Token (AVT).

- ⚡ QUIC trades performance for privacy [1]
- ⚡ Large certificates prolong handshakes [2]



QUIC Handshake

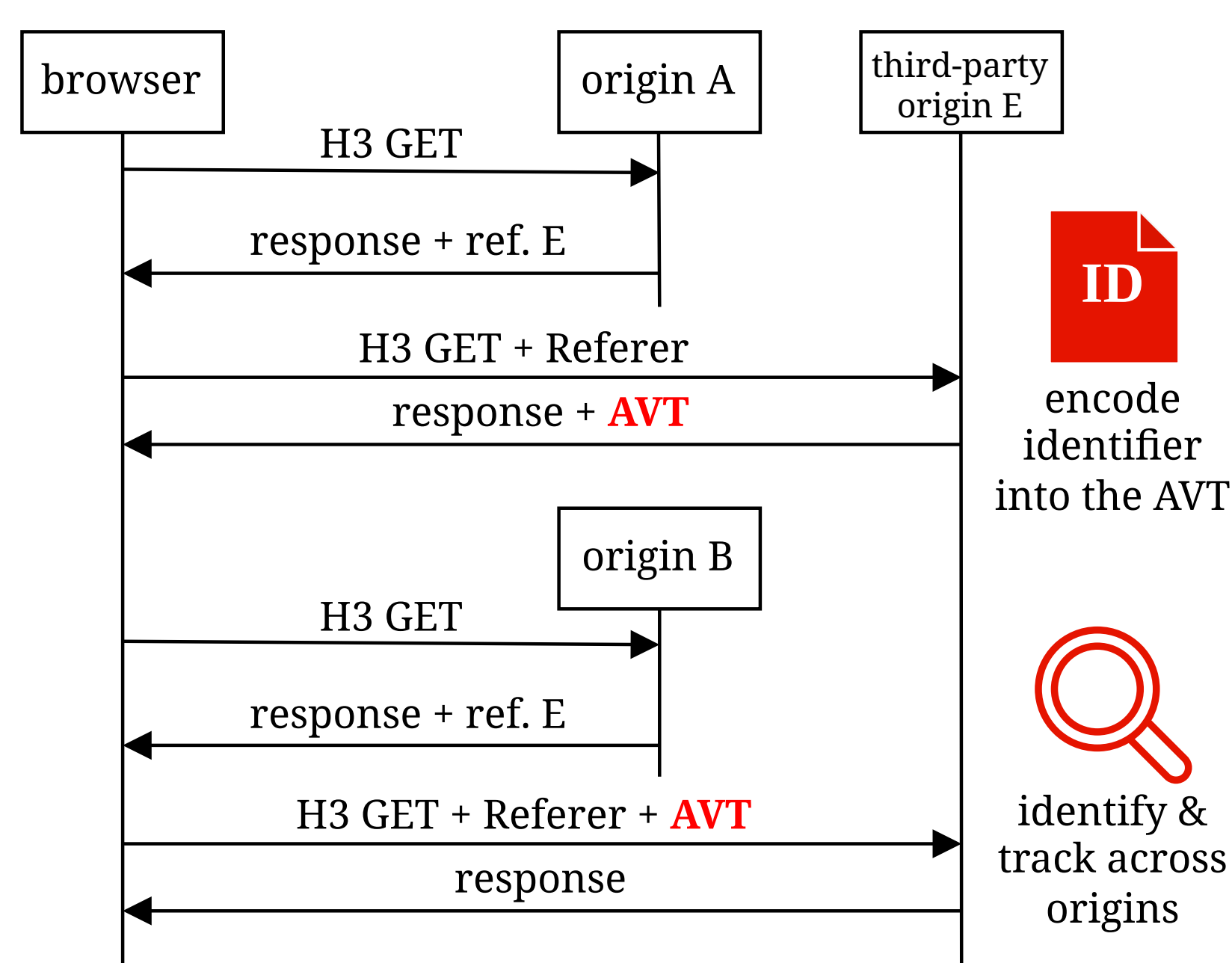


Many QUIC handshakes cannot be completed within 1 RTT because the response of the server exceeds the anti-amplification limit.

Address Validation Token (AVT)

- AVTs are shared by the QUIC server in *NEW_TOKEN* frames
- 15% of HTTP/3 servers issue AVTs, with sizes between 42 to 86 B
- All tested browsers cache AVTs for the whole browser session
- ⚡ AVTs are not available on first connect
- ⚡ Server can encode arbitrary data into an AVT

Web Tracking



AVTs can be used for tracking similar to cookies. Currently there is no evidence of AVT trackers.

Embedded third-parties identify origin by *Referer* header or origin specific URLs.

Web Tracking Protection

	Cookie Tracking	AVT Tracking
Chrome	❗ Vulnerable	❗ Vulnerable
Firefox	🛡 Protected	❗ Vulnerable
Brave	🛡 Protected	🛡 Protected

Firefox and Brave use isolated caches per origin to prevent tracking: *Total Cookie Protection* and *Ephemeral Storage*, respectively.

- ⚡ Common third-party page resources (e.g., fonts) benefit most from fast handshakes, but are also critical vectors for tracking

Privacy Pass (PP)

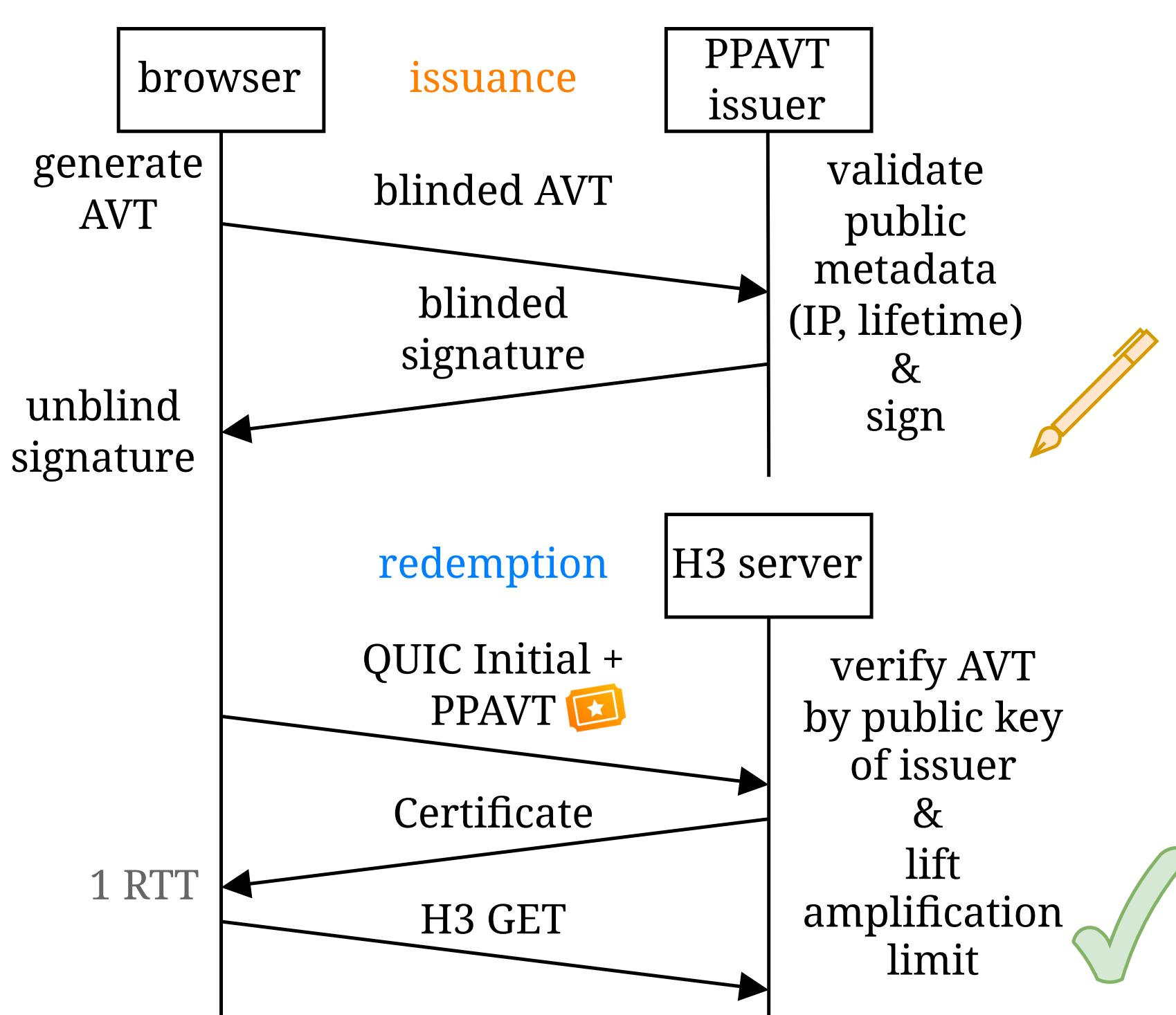
Privacy Pass [RFC 9576] enables privacy-preserving authentication.

Instead of presenting linkable state-carrying information to servers (e.g., cookies, AVTs), clients present unlinkable tokens, only sharing one-bit of information.

More information can be shared, as specified in the *public-metadata-issuance* draft, based on *PBRSA*.

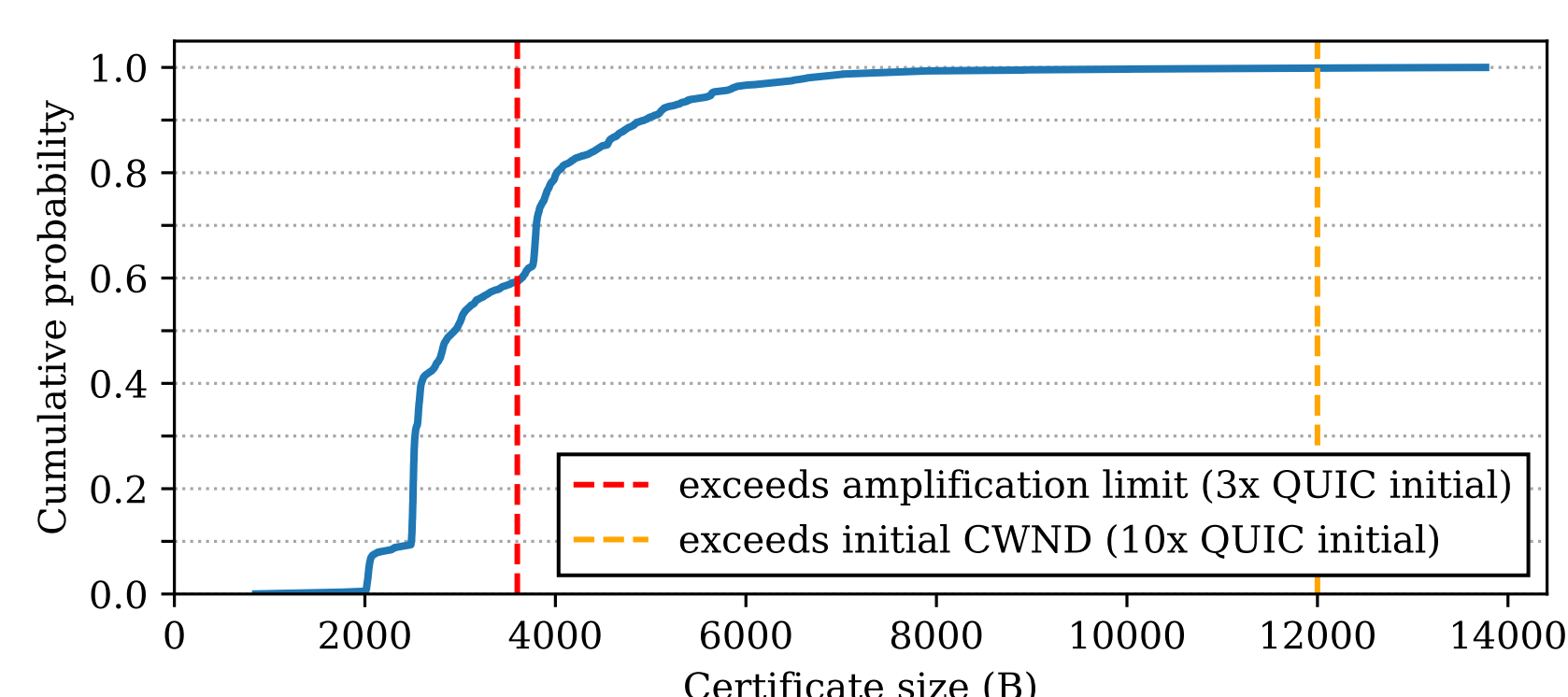


Privacy Pass Address Validation Tokens (PPAVT)



- ✓ Can lift anti-amplification limit on first connect
- ✓ Unlinkability of PP tokens prevents tracking

Why are first RTT responses long?



Distribution of certificate sizes in Tranco top 10k

- long certificate chains
- no coalesced QUIC
- no cert. compression
- no EC certificates

Cryptography (simplified)

Cryptography is based on *RSABSSA* [RFC 9474] and *draft-amjad-cfrg-partially-blind-rsa*.

1 The browser generates a partially blinded request using the issuer's public key pk_I .

```
nonce = random(32)
ext = { ip, lifetime }
blind_msg = Blind(pkI, nonce, ext)
req = { blind_msg, ext }
```

2 The issuer signs the request, after verifying the client's IP and lifetime.

```
blind_sig = BlindSign(skI, req)
resp = { blind_sig }
```

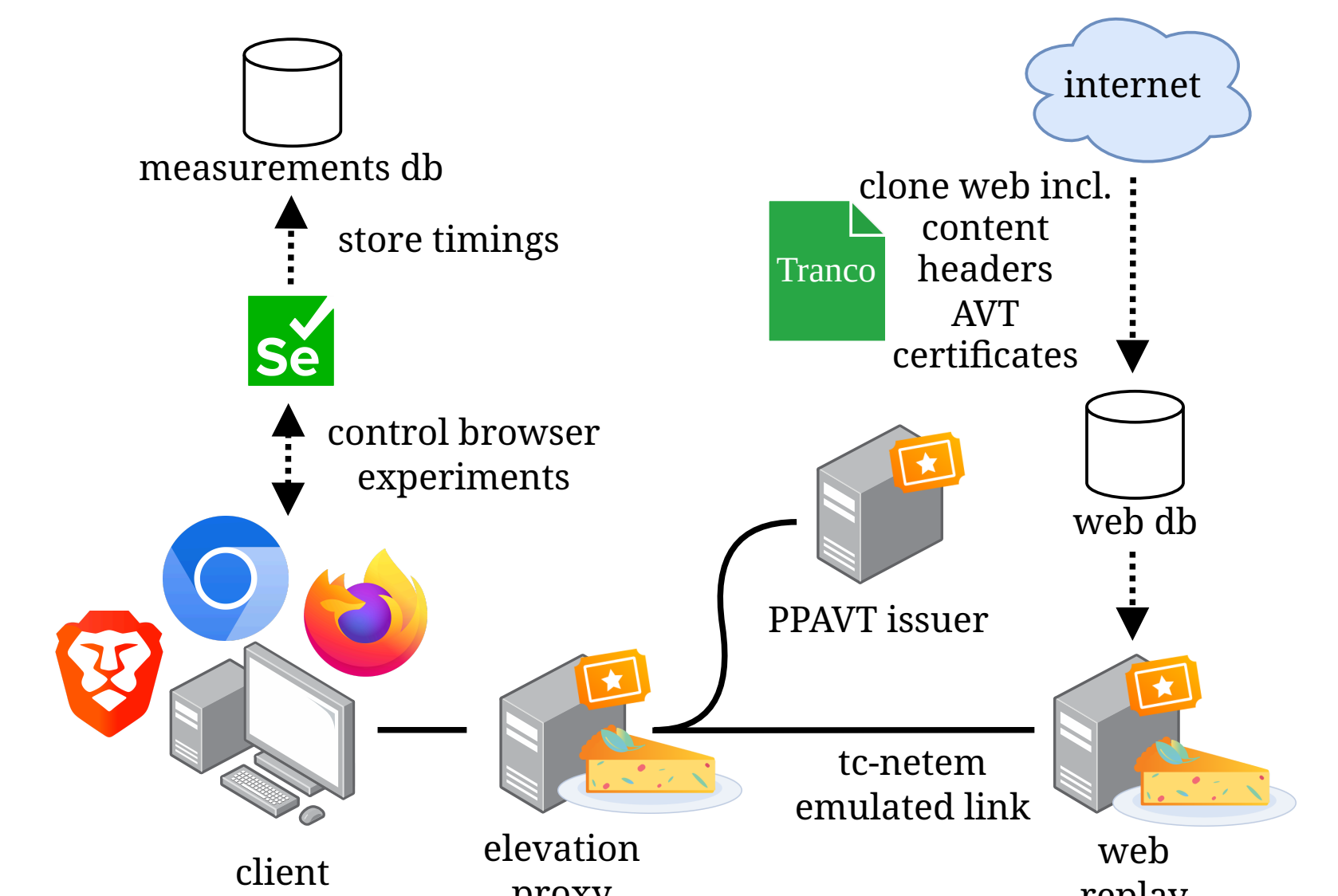
3 The client unblinds the signature, and generates the PPAVT.

```
sig = Finalize(pkI, nonce, ext, resp)
token = { nonce, ext, sig, issuer_id }
```

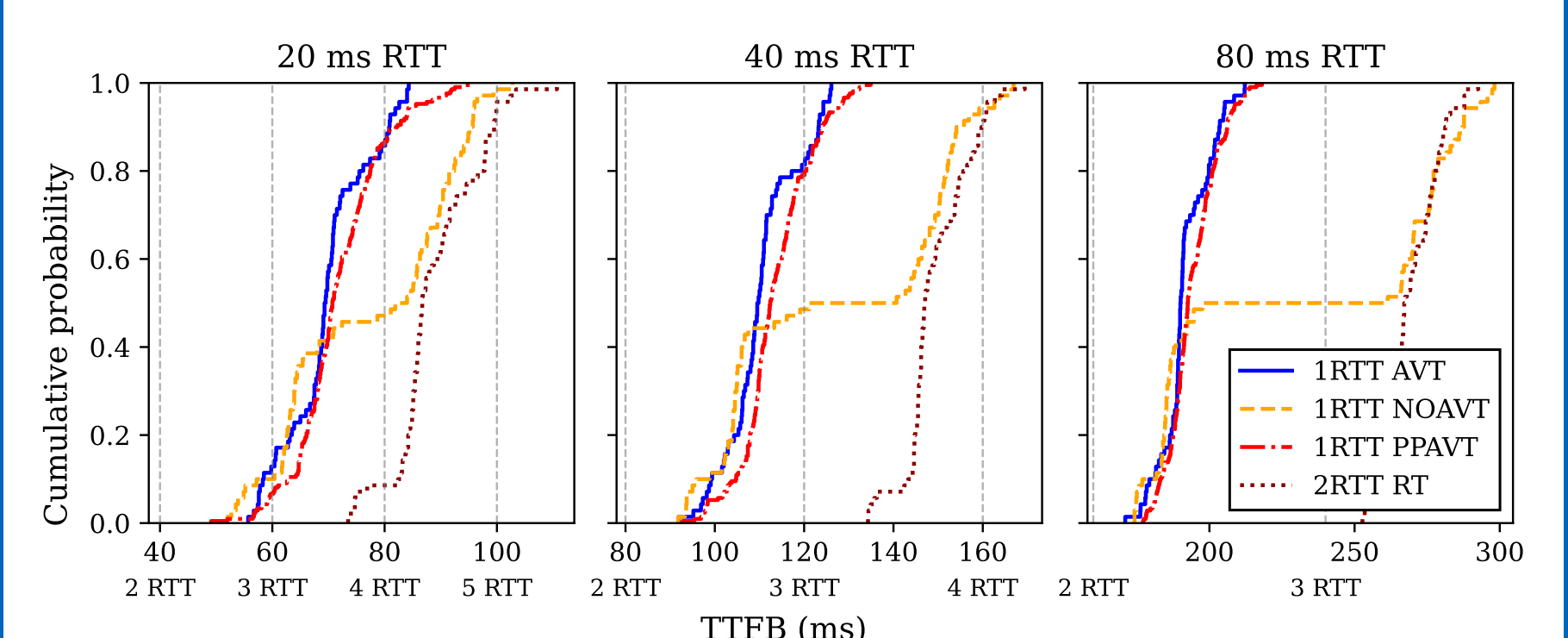
4 The server verifies the IP, lifetime, and signature with the issuer's public key.

```
Verify(pkI, token)
```

Testbed



Evaluation



Distribution of TTFB of different handshake modes

- ✓ PPAVT almost reduces handshake to server-provided AVT
- ⌚ Verification by Privacy Pass took about 9 ms

Open Challenges

- Address replay and double-spending problem
- Enhance browser and H3 origin replication
 - Evaluate the impact of *ML-KEM*
 - Integrate certificate compression
- Evaluate page load times
- Evaluate more than Tranco top 10
- Cover browsers beyond Chromium
- Optimize PP verification performance

[1] Erik Sy, Christian Burkert, Hannes Federrath, and Mathias Fischer. A QUIC Look at Web Tracking. In *PoPETs '19*, volume 2019, pages 255–266.

[2] Marcin Nawrocki, Pouyan Fotouhi Tehrani, Raphael Hiesgen, Jonas Mücke, Thomas C. Schmidt, and Matthias Wählisch. On the interplay between TLS certificates and QUIC performance. In *CoNEXT '22*, pages 204–213. ACM.